

Review

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Overview

- Stack Frames
- Stack Overflows
- ROP
- Mitigations
 - NX
 - ASLR
 - PIE
 - Stack Canaries
- Elf Structure
 - GOT+PLT
- Mitigation
 - RELRO

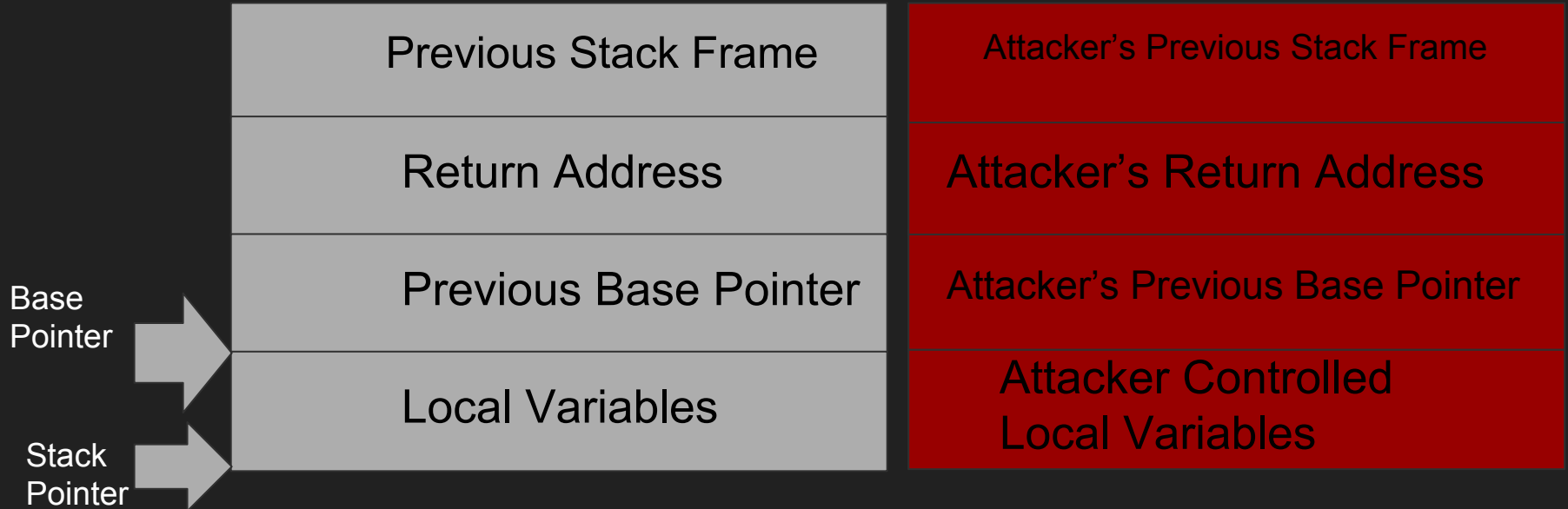
Stack Frames

- Contains information about function
 - Return address
 - Previous Base Pointer
 - Local Variables
 - In some cases parameters



Stack Overflows

- We read too much data onto the stack
 - Overwrite everything!



ROP

- Be intelligent in our overwriting
 - *ret == pop rip*
 - If we overwrite the return address to point to a ROP gadget, we can execute multiple pieces of code
 - ROP gadget: something ending with *ret*

Attacker's #n Gadget Address

Attacker's #3 Gadget Address

Attacker's #2 Gadget Address

Attacker's #1 Gadget Address

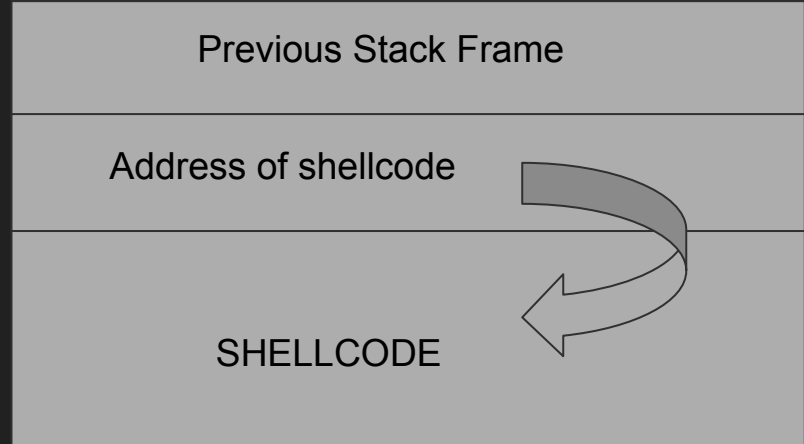
Attacker's Previous Base Pointer

Attacker Controlled
Local Variables

Stack Overflows - Shellcode

- Older way of exploiting overflows
- Point the return address into a buffer we control
 - Contains “shellcode”
 - Code that gives us a shell when run

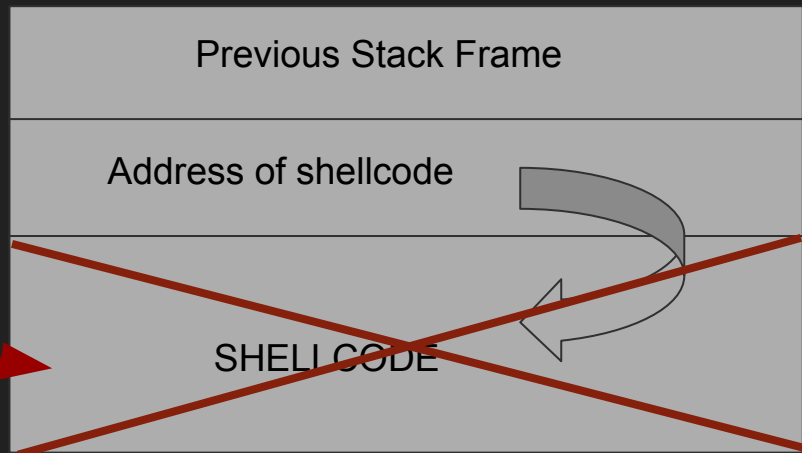
Shellcode itself can sometimes be used, but this method is outdated



Mitigations - NX

- NX / DEP / W^X
 - Non Executable
 - Data Execution Prevention
 - Writable XOR Executable
- What does this mean for us?
 - The stack cannot be executed
 - No shellcode
 - D:
- How to bypass?
 - ROP

Cannot be
executed



Mitigations - ASLR

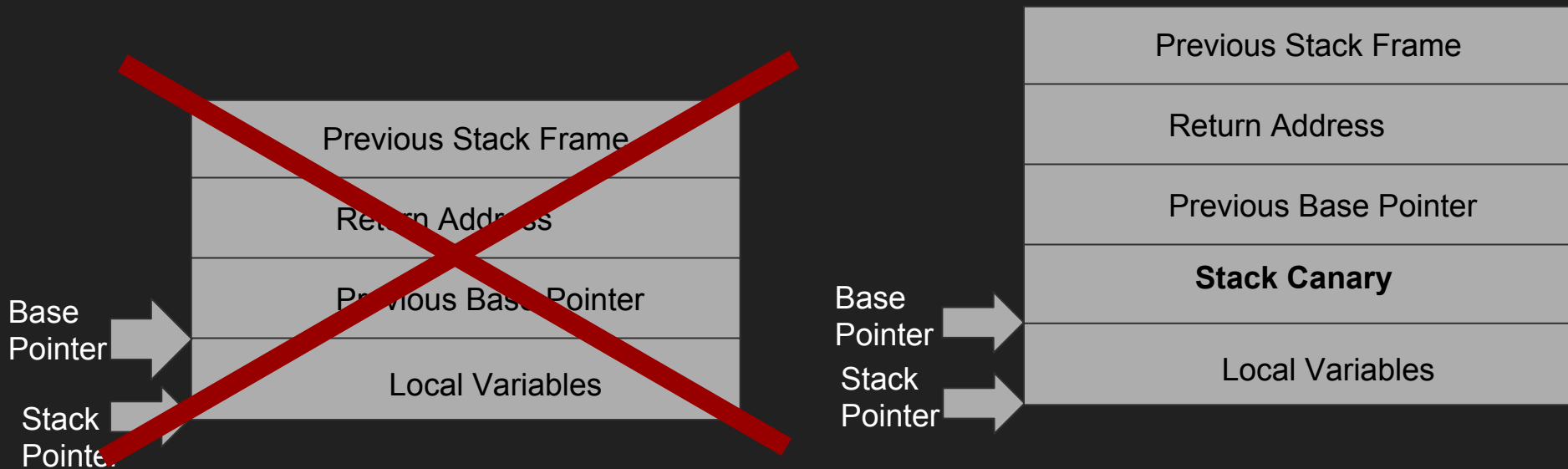
- Address Space Layout Randomization
- You can't return to what you can't find!
 - Randomizes libc
 - Randomizes Stack
 - Randomizes Heap
- Notably, our executable does NOT get randomized
- How to bypass
 - Leak Pointer
 - Partial Overwrite #we haven't talked about
 - Ret2plt #we haven't talked about
 - ROP in target executable

Mitigations - PIE

- Position Independent Execution
- ASLR -- But actually fully applied
- Randomizes the executable location
- Technically within the subset of ASLR
- How to bypass:
 - Leak Pointer
 - Partial Overwrite #we haven't talked about

Mitigations - Stack Canaries

- Make stack overflows impossible to exploit
- Put random value on stack
- Check that it hasn't changed before returning



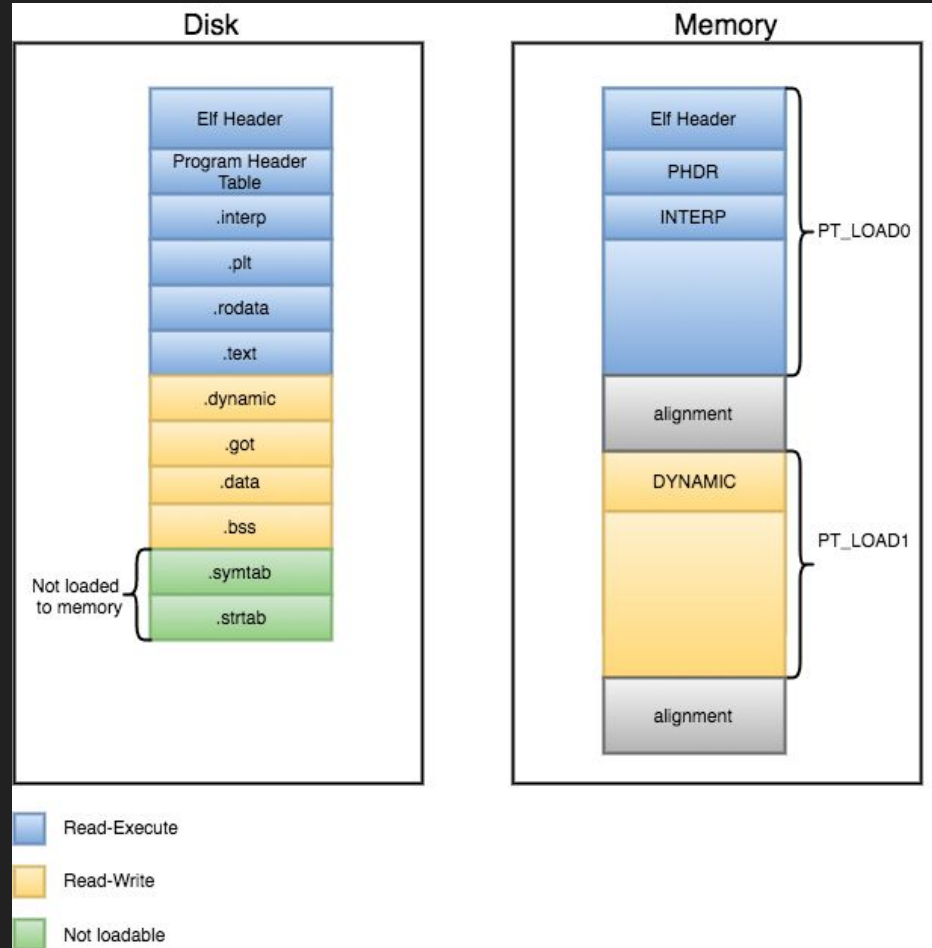
Tool - Checksec

- How to tell which ones are enabled
 - `checksec /path/to/file`
 - Comes with `pwntools`
- ASLR is actually enabled system wide, it's by default on

```
[st@localhost got1]$ checksec ./got1
[*] '/home/st/Desktop/teaching/3/got1/got1'
Arch:          amd64-64-little
RELRO:         Partial RELRO
Stack:         No canary found
NX:            NX enabled
PIE:           No PIE (0x400000)
```

Elf Structure

- Lots of sections
- Interesting Ones:
 - .text
 - .plt
 - .got
- No Write+Execute
 - That's DEP at work!



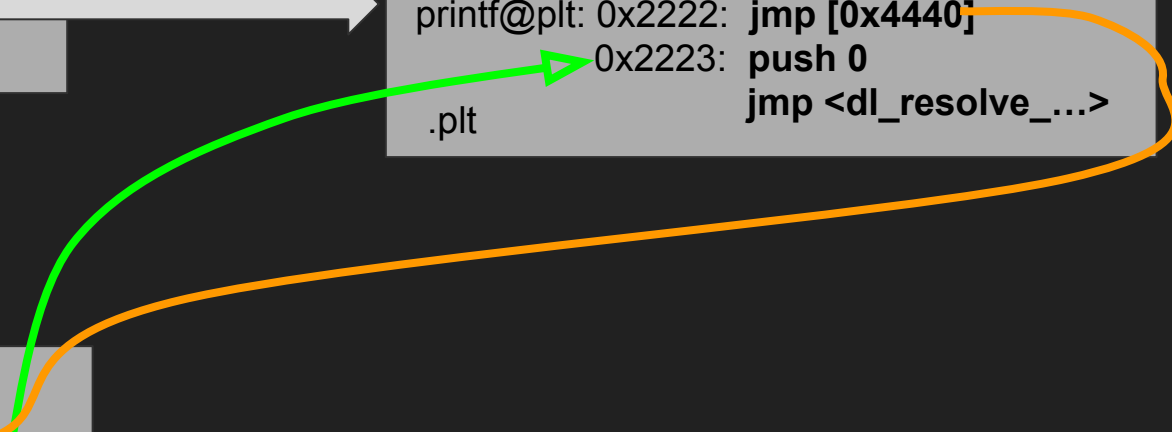
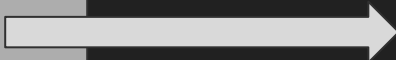
Sections

- `.text`
 - Contains all the code we write
 - Just assembly
- `.plt`
 - Trampoline for external calls
 - Example: `printf`, `puts`, `fgets`, `system`, ...
 - Handles dealing with ASLR to find true addresses of external functions
- `.got`
 - Table of addresses
 - Used by `plt` to store true addresses of external functions

```
mov rdi, 0x20131
call printf@plt
.text
```

```
printf@plt: 0x2222: jmp [0x4440]
0x2223: push 0
jmp <dl_resolve_...>
.plt
```

```
Before First Call
printf@got: 0x4440: 0x2223
puts@got: 0x4448: 0x2226
.got
```



```
mov rdi, 0x20011
call printf@plt
mov rdi, 0x20131
call printf@plt
```

.text

```
printf@plt: 0x2222: jmp [0x4440]
0x2223: push 0
jmp <dl_resolve_...>
```

.plt

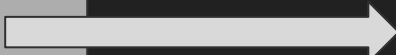
After Being Called

```
printf@got: 0x4440: 0x9999
puts@got: 0x4448: 0x2226
```

.got

```
printf: 0x999: push rbp
mov rbp, rsp
...
ret
```

.text of libc



Mitigation - RELRO

- RELOCations Read Only
- Because of how the GOT is lazily loaded, it is writeable
- Full RELRO:
 - GOT is filled at load time rather than runtime
 - GOT is not writeable
- Partial RELRO:
 - Default
 - Doesn't mean anything to us

Questions or Comments

- Anything anyone is remotely hazy on